

Min Mux
 $Q_7: SOP / POS$

1) Start for using bool alg / K-map:

$F(A, B, C) = \sum(0, 1, 4, 6)$
 $F(A, B, C) = \sum(m_0, m_1, m_4, m_6)$

$f = \overline{A}\overline{B}\overline{C} + \overline{A}B\overline{C} + \overline{A}B\overline{C} + \overline{A}B\overline{C}$
 $f = \overline{A}\overline{B} + \overline{A}C$
SOP

POS
 $Q = \overline{10}(0, 1, 4, 8, 9, 12, 15)$

$f = \overline{C}\overline{D} + \overline{A}C + ABCD$
 $f = (\overline{C}\overline{D})(\overline{A}C)(ABCD)$
 $f = (\overline{C} + \overline{D})(\overline{B} + \overline{C})(\overline{A} + B + \overline{C} + D)$
 $f = (C + D)(B + C)(\overline{A} + \overline{B} + \overline{C} + D)$

Q2: Funct PUN / PON
 1) CMOS = NMOS + PMOS

NAND

AND

NOR

OR

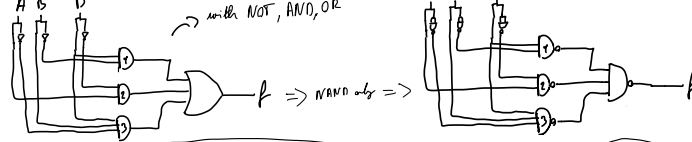
| x_1 | x_2 | x_3 | x_4 | NAND | AND | NOR | OR |
|-------|-------|-------|-------|------|-----|-----|----|
| 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 0 | 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 |
| 0 | 0 | 1 | 1 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 1 | 1 | 0 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 | 1 | 0 |
| 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 |
| 1 | 0 | 1 | 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 | 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 0 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 1 | 1 | 0 | 1 | 0 |

$f \rightarrow CMOS: + \rightarrow parallel, \rightarrow series, \#N = \#P$

Bool algebra:

- $x \cdot 0 = 0, x \cdot 1 = x, x \cdot x = x, x \cdot \overline{x} = 0, x \cdot \overline{\overline{x}} = x$
- $x + 0 = x, x + 1 = 1, x + x = x, x + \overline{x} = 1, x + \overline{\overline{x}} = x$
- $x \cdot (y + z) = x \cdot y + x \cdot z$ (Distribution)
- $x + x \cdot y = x$ (Absorption)
- $x \cdot y + x \cdot \overline{y} = x$ (Combining)
- $\overline{x \cdot y} = \overline{x} + \overline{y}$ (DeMorgan's theorem)
- $x + \overline{\overline{x}} = x + x = x$ (Consensus)

2) Funct \rightarrow NAND / NOR gate pos: \rightarrow with NOT, AND, OR

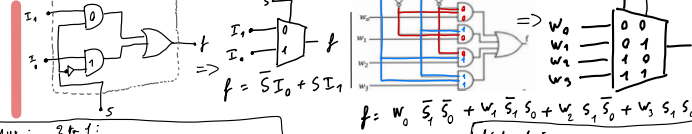


$y = \overline{A\overline{B} + \overline{B}CD} = \overline{\overline{B}(A + CD)}$
 $y = \overline{\overline{B}} \cdot \overline{(A + CD)} = B \cdot (\overline{A} + \overline{C} + \overline{D})$
 $y = \overline{B} + \overline{A} + \overline{C} + \overline{D}$
 $y = \overline{B} + \overline{A} \cdot \overline{C} \cdot \overline{D}$
 $y = \overline{B} + \overline{A} \cdot (\overline{C} + \overline{D})$

Half Adder:

| x | y | Sum | Carry |
|-----|-----|-----|-------|
| 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |

3) Funct using Multiplexers (MUX):



MUX: 2 to 1:
 $f = w_0 w_2 + w_1 w_3 + w_2 w_3$
 $f = \overline{w_1} (0 \cdot w_2 + 0 \cdot w_3 + w_2 w_3) + w_1 (1 \cdot w_2 + 1 \cdot w_3 + w_2 w_3)$
 $f = \overline{w_1} (w_2 w_3) + w_1 (w_2 + w_3 + w_2 w_3)$
 $f = \overline{w_1} (w_2 w_3) + w_1 (w_2 + w_3)$
 $f = \overline{w_1} g + w_1 h \rightarrow$ expand by w_2
 $g = w_2 w_3, h = w_2 + w_3$
 $g = \overline{w_2} (0 \cdot w_3) + w_2 (1 \cdot w_3) = \overline{w_2} (0) + w_2 (w_3)$

4 to 1:
 $f = w_0 w_2 + w_1 w_3 + w_2 w_3$
 $f = \overline{w_1} \overline{w_2} (0 + 0 + 0) + \overline{w_1} w_2 (0 + 0 + w_3) + w_1 \overline{w_2} (0 + w_3 + 0) + w_1 w_2 (1 + w_3 + w_3)$
 $= \overline{w_1} \overline{w_2} (0) + \overline{w_1} w_2 (w_3) + w_1 \overline{w_2} (w_3) + w_1 w_2 (1)$

Full adder:

| C_{in} | I_1 | I_2 | C_{out} | S |
|----------|-------|-------|-----------|-----|
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 1 | 1 |

4 Bit adder:

4) DeMorgan's theorem:

$\overline{x \cdot y} = \overline{x} + \overline{y}$
 $\overline{x + y} = \overline{x} \cdot \overline{y}$

SOP \rightarrow **POS**
 $\overline{w_1} \rightarrow 0, \overline{w_1} \rightarrow 1$

2 to 1 mux \rightarrow 1 Variable
4 to 1 mux \rightarrow 2 Variables

most significant
least significant

Overflow: Sign bit is wrong for the equation

- if the result is $> 2^{N-1} - 1$
- if the result is $< -2^{N-1}$

Signal

Unsigned:
 $0 \rightarrow 255$
 $2^8 \cdot 2^3 + 2^2 = 139_{10}$

8-bit:
 $-128 \rightarrow 127$
 $2^6 \cdot 2^3 + 2^2 = 110_{10}$
 $2^5 \cdot 2^3 + 2^2 = 110_{10}$
 $-2^3 + 2^2 + 2^2 + 2^2 = -4_{10}$